
SUPPLEMENTAL MATERIALS FOR: BEYOND DATA AND MODEL PARALLELISM FOR DEEP NEURAL NETWORKS

Algorithm 1 Full Simulation Algorithm.

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1: Input: An operator graph  $\mathcal{G}$ , a device topology  $\mathcal{D}$ , and a
  parallelization strategy  $\mathcal{S}$ .
2:  $\mathcal{T} = \text{BUILDTASKGRAPH}(\mathcal{G}, \mathcal{D}, \mathcal{S})$ 
3: readyQueue =  $\{\}$  // a priority queue sorted by readyTime
4: for  $t \in \mathcal{T}_N$  do
5:   t.state = NOTREADY
6:   if  $\mathcal{I}(t) = \{\}$  then
7:     t.state = READY
8:     readyQueue.enqueue(t)
9:   end if
10: end for
11: while readyQueue  $\neq \{\}$  do
12:   Task t = readyQueue.dequeue()
13:   Device d = t.device
14:   t.state = COMPLETE
15:   t.startTime =  $\max\{t.\text{readyTime}, d.\text{last.endTime}\}$ 
16:   t.endTime = t.startTime + t.exeTime
17:   d.last = t
18:   for  $n \in \mathcal{O}(t)$  do
19:     n.readyTime =  $\max\{n.\text{readyTime}, t.\text{endTime}\}$ 
20:     if all tasks in  $\mathcal{I}(n)$  are COMPLETE then
21:       n.state = READY
22:       readyQueue.enqueue(n)
23:     end if
24:   end for
25: end while
26: return  $\max\{t.\text{endTime} \mid t \in \mathcal{T}_N\}$ 

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1 FULL SIMULATION ALGORITHM

Algorithm 1 shows the pseudocode of the full simulation algorithm. It first builds a task graph using the method described in Section 5.1 and then sets the properties for each task using a variant of Dijkstra’s shortest-path algorithm (Cormen et al., 2009). Tasks are enqueued into a global priority queue when ready (i.e., all predecessor tasks are completed) and are dequeued in increasing order by their readyTime. Therefore, when a task t is dequeued, all tasks with an earlier readyTime have been scheduled, and we can set the properties for task t while maintaining the FIFO scheduling order (assumption A3).

2 DELTA SIMULATION ALGORITHM

Algorithm 2 shows the pseudocode of the full simulation algorithm. It first updates tasks and dependencies from an existing task graph and enqueues all modified tasks into

Algorithm 2 Delta Simulation Algorithm.

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1: Input: An operator graph  $\mathcal{G}$ , a device topology  $\mathcal{D}$ , an original
  task graph  $\mathcal{T}$ , and a new configuration  $c'_i$  for operator  $o_i$ .
2: updateQueue =  $\{\}$  // a priority queue sorted by readyTime
3: /*UPDATETASKGRAPH returns the updated task graph and a
  list of tasks with new readyTime*/
4:  $\mathcal{T}, \mathcal{L} = \text{UPDATETASKGRAPH}(\mathcal{T}, \mathcal{G}, \mathcal{D}, c_i, c'_i)$ 
5: updateQueue.enqueue( $\mathcal{L}$ )
6: while updateQueue  $\neq \{\}$  do
7:   Task t = updateQueue.dequeue()
8:   t.startTime =  $\max\{t.\text{readyTime}, t.\text{preTask.endTime}\}$ 
9:   t.endTime = t.startTime + t.exeTime
10:  for  $n \in \mathcal{O}(t)$  do
11:    if UPDATETASK( $n$ ) then
12:      updateQueue.push( $n$ )
13:    end if
14:  end for
15:  if UPDATETASK( $t.\text{nextTask}$ ) then
16:    updateQueue.push( $t.\text{nextTask}$ )
17:  end if
18: end while
19: return  $\max\{t.\text{endTime} \mid t \in \mathcal{T}_N\}$ 
20:
21: function UPDATETASK( $t$ )
22:   t.readyTime =  $\max\{p.\text{endTime} \mid p \in \mathcal{I}(t)\}$ 
23:   /*Swap  $t$  with other tasks on the device to maintain
  FIFO.*/
24:   t.startTime =  $\max\{t.\text{readyTime}, t.\text{preTask.endTime}\}$ 
25:   if  $t$ ’s readyTime or startTime is changed then
26:     return True
27:   else
28:     return False
29:   end if
30: end function

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a global priority queue (line 4-5). Similar to the Bellman-Ford shortest-path algorithm (Cormen et al., 2009), the delta simulation algorithm iteratively dequeues updated tasks and propagates the updates to subsequent tasks (line 6-14). The full and delta simulation algorithms always produce the same timeline for a given task graph.

REFERENCES

Cormen, T. H., Leiserson, C. E., Rivest, R. L., and Stein, C. *Introduction to Algorithms, Third Edition*. The MIT Press, 3rd edition, 2009.